



Sanitizable Signatures in XML Signature

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Sanitizable Signatures in XML Signature — Performance, Mixing Properties, and Revisiting the Property of Transparency

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A Sanitizable Signature Scheme allows

- a defined third party („sanitizer“)
- to alter defined parts of an already signed document
- without invalidating the given signature,
- without interaction between signer and sanitizer.

Sanitizable Signature Scheme

Generally consists of five algorithms:

1. **Setup**: Generate key pair and public parameters
2. **Sign**: Generate a Signature over sanitizable and immutable parts
3. **I-Forge**: Change sanitizable parts such that signature is still valid.
Requires knowledge of the „sanitizer secret“.
4. **U-Forge**: Change sanitizable parts such that signature is still valid.
Requires two different „versions“, i.e. original and sanitized msg.
5. **Verify**: Verify signature's validity. Valid iff immutable parts are unchanged and changing the sanitized parts using I-Forge or U-Forge.

Contribution

- **Our Implementation of 5 Schemes in JAVA yields:**
 - Tolerable Performance penalty compared to SHA/RSA (most of the schemes)
 - Integration into JAVA Crypto Framework (JCA) possible (as a new JAVA Crypto Provider)
 - Integration into XML Digital Signature Syntax and Processing Standard (W3C) possible
- **Changing Properties by Mixing Chameleon- and SHA-Hashes**
- **More Precise Definition of 3 Properties Transparency and Strong- / Weak-Transparency**

Research Context



Project Goal: IT-supported Robust & Secure Supply Chains

Goal of University of Passau / Institute of IT Security and Security Law:

- **Integrity and Authenticity Statements for Partial Data**
- **Legal Compliance, Manage & Verify the Statements**

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**Federal Ministry
of Education
and Research**

Why Sanitizable Signatures

- Notion appeared in literature around 2005 in work of G. Ateniese, D. H. Chou, B. de Medeiros, and G. Tsudik.
- Concept also described in earlier works, i.e. Content Extraction Signatures by Steinfeld et al. in 2001
- Ever since: Many Schemes
- Many Use Cases:
i.e. Freight Document
- No Implementations
- Not applied on a large scale

Super Glue Producer GmbH & Co KG



Telefon 123349 Telefax 123348

RoadSupply
2nd Left Street
PX93774 Big Town

381377 LIEFERSCHEIN

Nummer :0000 16 v26.01.2011
Kunde :ID 73
Auftrag :10000136491 v26.01.2011
Bestellnr.:

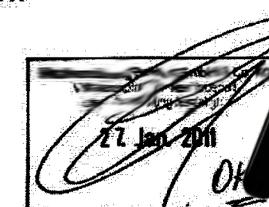
Pos.	Artikelnummer Bezeichnung	ME	Liefer-
1	SUP12 Superglue: SUPERIOR TW	kg	350
2	12 Kanister 11 KANISTER	kg	20
3	SUP007 Superglue: SUPERIOR VPE	kg	330
4	4 Kanister	kg	

Zustand des Fahrzeugs incl. Ladungssicherung wurde vor der Ver-
überprüft!

Der abholende Fahrer verpflichtet sich durch seine Unterschrift, die
Ladungssicherungsmittel (Zurrgurte, Antirutschmatten, Sperrstangen etc.)
für alle Packstücke bei Verladung eingesetzt zu haben und auf eine
verpflichtende Kontrolle der Ladungssicherung nach max. 50 km Fahrtstrecke
hingewiesen worden zu sein.

Anlieferung bis 28. 1.

Achtung: Thermoware > 10 °C!!!!



Ordnung gemäß Temperiert und Transportiert!
AMTL. KNZ. KA AB 285

Bei Lieferung 1 Kanistr besch.

2701 / MSS Pauschalent



Five Sanitizable Signature Schemes

Ateniese et al.'s Scheme:

- Sanitizable Signature based on Chameleon hashes (CH)
- Sanitizer can compute hash collisions, if trapdoor information is known
- Different CH usable within Ateniese scheme :
 1. **Krawczyk**: 1st chameleon hash, based on DLP assumption
 2. **Ateniese**: ID-based approach
 3. **Zhang**: ID-based approach without an UForge-algorithm
 4. **Chen**: ID-based approach without the key-exposure-problem
- 5. **Miyazaki et al.'s scheme**:
 - redactable signature scheme based on commitments
 - allows just deletion
 - controllable redaction of consecutive sanitizers

XML Signature Syntax and Processing W3C Standard

```
1  <?xml version="1.0" encoding="UTF-8" standalone="no"?>
2  <PurchaseOrder>
3    <Item id="8492341">
4      <Description id="8492340">Video Game</Description>
5      <Price>10.29</Price>
6    </Item>
7    <Signature xmlns="http://www.w3.org/2000/09/xmldsig#">
8      <SignedInfo>
9        <CanonicalizationMethod Algorithm="http://www.w3.org/2001/10/xml-exc-c14n#" />
10       <SignatureMethod Algorithm="http://www.w3.org/2000/09/xmldsig#rsa-sha1"/>
11       <Reference URI="#xpointer(id('8492340'))">
12         <Transforms>
13           <Transform Algorithm="http://www.w3.org/2000/09/xmldsig#enveloped-signature"/>
14         </Transforms>
15         <DigestMethod Algorithm="http://www.w3.org/2000/09/xmldsig#sha1"/>
16         <DigestValue>ABYPTWCzr8F7dx1UKyg1C+tycm4=</DigestValue>
17       </Reference>
18     </SignedInfo>
19     <SignatureValue>D9hok43bgiRJ9uzp/7A9MA2YZBFuivvzoZTbC(....)DsFCXjtkRxQ==</SignatureValue>
20   </Signature>
21 </PurchaseOrder>
```

The diagram illustrates the XML signature structure with numbered steps:

- Step 1: A callout points to the `<Reference URI="#xpointer(id('8492340'))">` element.
- Step 2: An arrow points from the `URI="#xpointer(id('8492340'))"` attribute to the `<Transforms>` element.
- Step 3: An arrow points from the `<Transforms>` element to the `<DigestValue>` element.
- Step 4: An arrow points from the `<DigestValue>` element back to the `<Reference URI="#xpointer(id('8492340'))">` element.

Sanitizable XML Signature (W3C standard compliant)

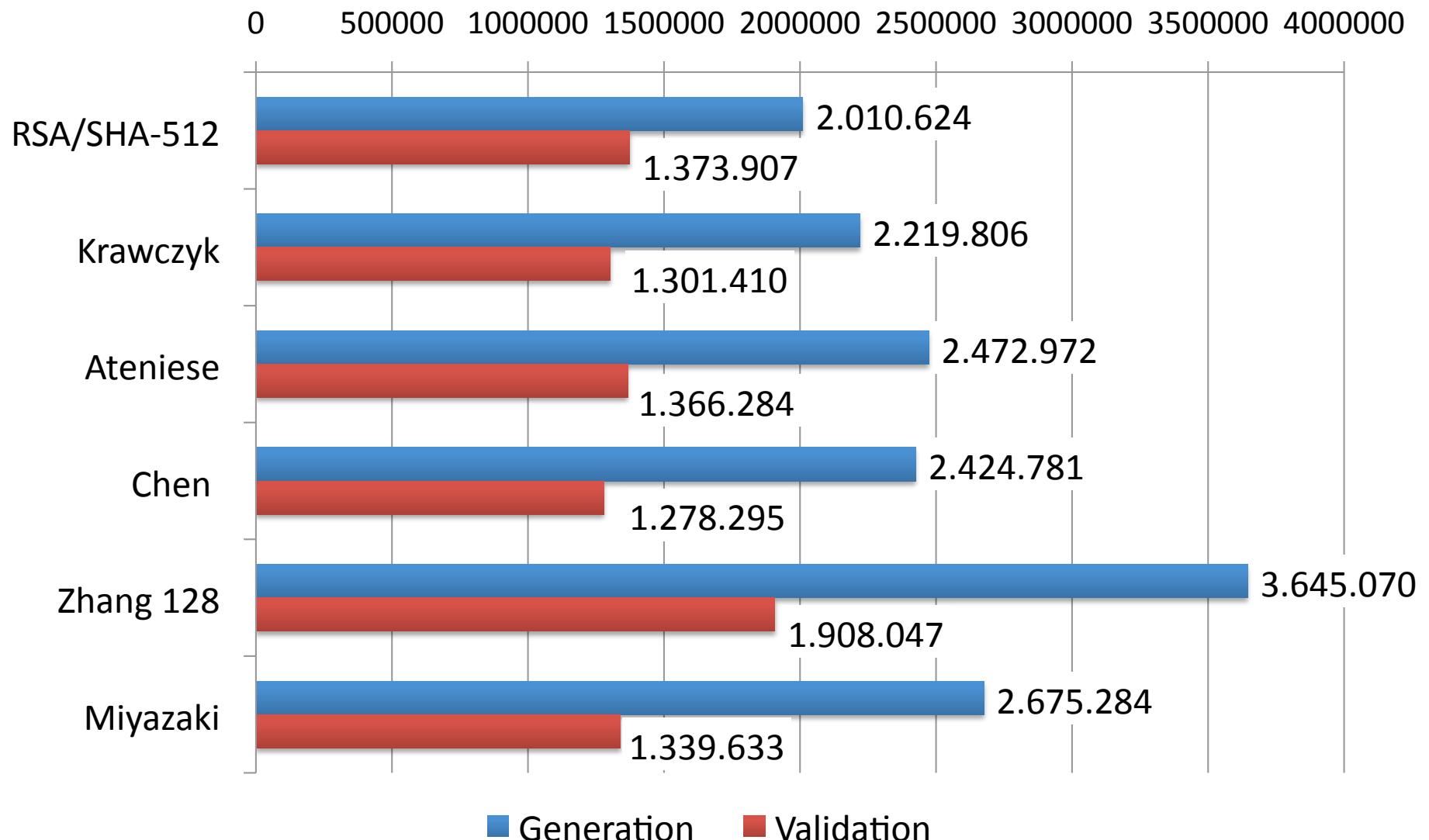
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1  <?xml version="1.0" encoding="UTF-8" standalone="no"?>
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11       <Reference URI="#xpointer(id('8492340'))"> 2.
12         <Transforms>
13           <Transform Algorithm="http://www.w3.org/2000/09/xmldsig#enveloped-signature"/>
14         </Transforms>
15         <DigestMethod Algorithm="http://www.example.org/xmldsig-more#chamhashdisc"> 3.
16           <ChamHashDiscKeyValue>
17             <p>Aa5Mue7ppx2YD7R8KXUq0IKSTSay6jHhWm9L0dxHpL2P</p>
18             <q>1yZc93TTjswH2j4UupUgQUKmk11GPCtN6Xo7iPSXsc=</q>
19             <r>FQrJPkWb0JwiffjrAdbWAoyropQmNohMgEy6ABsvptQ=</r>
20             <g>JtqJ1H0NL0Is+6Y797XKQ1hbHc+HYgoGQAkvK8h+q8Y=</g>
21             <y>AVwdxM1XF6HIHRH10r7Xoojb0VoB7ZBP4Dxc83BDDgxG</y>
22           </ChamHashDiscKeyValue>
23         </DigestMethod>
24         <DigestValue>8Xt2AtyvB3Umwf8LlyrGSVnvLc4=</DigestValue> 4.
25       </Reference>
26     </SignedInfo>
27     <SignatureValue>D9hok43bgiRJ9uzp/7A9MA2YZBFuivvzoZTbC(....)DsFCXjtkRxQ==</SignatureValue>
28   </Signature>
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Sanitizable XML Signature (W3C standard compliant)

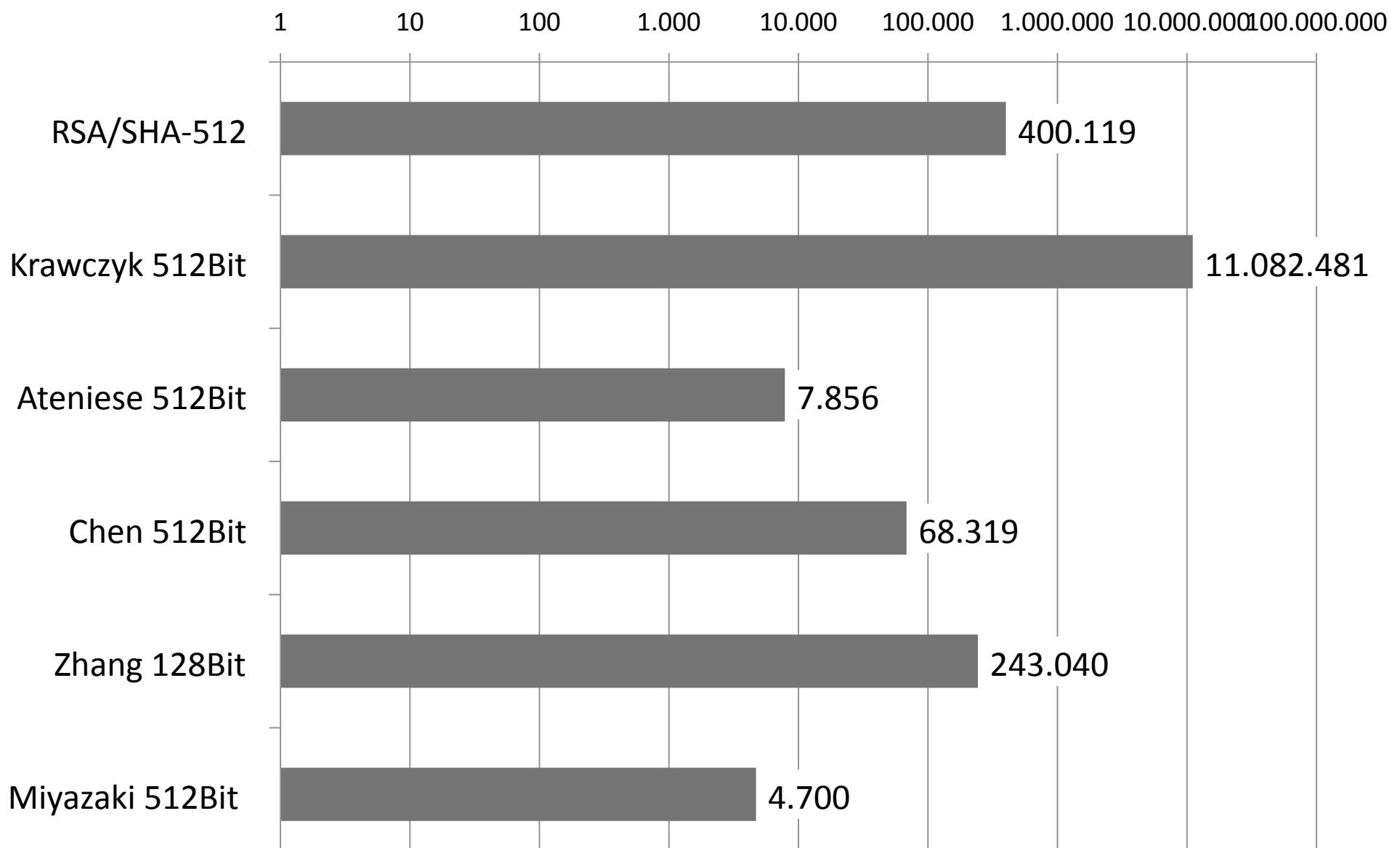
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22       <p>Aa5Mue7ppx2YD7R8KXUqQIKSTSay6jHhWm9L0dxHpL2P</p>
23       <q>1yZc93TTjswH2j4UupUgQUkmk111GPCtN6Xo7iPSXsc=</q>
24       <r>FQrJPkWb0JwiffjrAdbWAoyropQmNohMgEy6ABsvptQ=</r>
25       <g>JtqJ1H0NL0Is+6Y797XKQ1hbHc+HYgoGQAkvK8h+q8Y=</g>
26       <y>AVwdxM1XF6HIHRH10r7Xoojb0VoB7ZBP4Dxc83BDDgxG</y>
27     </ChamHashDiscKeyValue>
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```

The diagram illustrates the sanitization process of an XML signature. It shows the original XML code with various parts highlighted by orange arrows and circles. Step 1 highlights the entire XML structure. Step 2 highlights the reference to the signed item. Step 3 highlights the digest method. Step 4 highlights the digest value. A 'new' label points to the ChamHashDiscKeyValue element.

Performance: Sign & Verify (No Setup) in micro secs.



Performance: Setup incl. KeyPair Generation in micro secs.



Performance: Summary

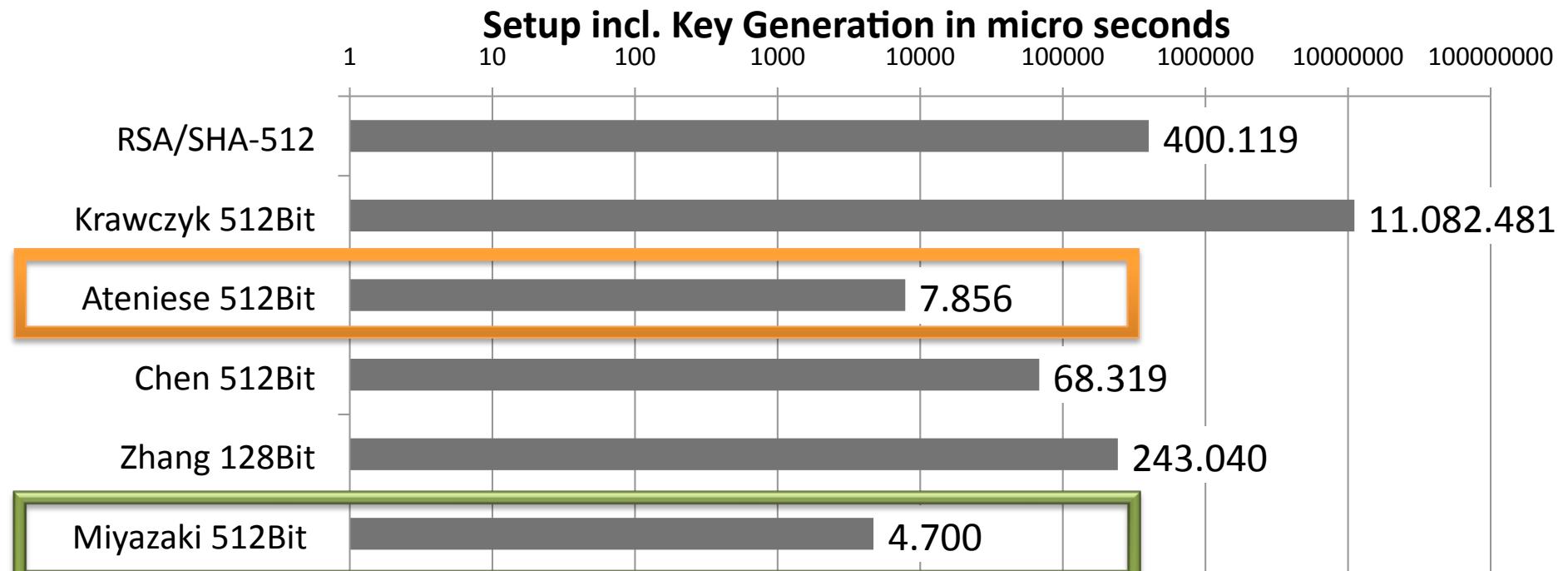
excluding key generation & setup:

similar runtime as SHA/RSA for signature generation and validation

exception: Zhang et al.'s scheme based on Elliptic Curve Crypto

including key generation & setup:

Overall: Chameleon Hash by Ateniese et al. performs best



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Overall: Chameleon Hash by Ateniese et al. performs best

comparison of just one execution is not always enough:

key exposure problem

**“a forged message and a original message leaks the secret
and allows to U-Forge other messages under the same key”**

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excluding key generation & setup:

similar runtime as SHA/RSA for signature generation and validation

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Overall: Chameleon Hash by Ateniese et al. performs best

comparison of just one execution is not always enough:

suffer from key exposure

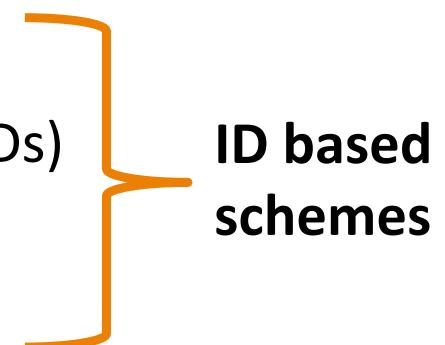
Krawczyk et al. (needs new key for each message)

key exposure problem reduction possible

Ateniese et al. (using one time TransactionIDs)

key exposure free

Chen et al. & Zhang et al.



Example: Ateniese Scheme for Chameleon Hashes

Message is split into parts (m_j) , each hashed independently

Message: $\textcolor{brown}{m} = m_1 \parallel \dots \parallel m_i$

\mathcal{CH} : a Chameleon Hash function

k_i : secret key given to the Sanitizer needed for I-Forge

u_i : secret key that is safely discarded

Signature:

$$d_i = \begin{cases} \mathcal{CH}_{k_i}(m_i) & \text{if } m_i \text{ is sanitizable} \\ \mathcal{CH}_{u_i}(m_i) & \text{if } m_i \text{ is not sanitizable} \end{cases}$$

$$\sigma = SIGN(d_1 \parallel \dots \parallel d_n)$$

Transparency, Weak Transparency, Strong Transparency

Ateniese et al. define “Transparency”:

*Given a signed message with a valid signature,
no party – except the signer and the censor –
should be able to correctly guess whether
the message has been sanitized.*

Ateniese et al. divided this into

“Weak Transparency”:

*(...) the verifier knows exactly which parts of the message
are potentially sanitizable (...)*

“Strong Transparency”:

*(...) the verifier does not know which parts of the message (...)
could potentially be sanitizable.*

Properties of Ateniese Scheme for Chameleon Hashes

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Properties: TRANSPARENCY & STRONG TRANSPARENCY

Mixing

Message is split into parts (m_j) , each hashed independently

Message: $\textcolor{brown}{m} = m_1 \parallel \dots \parallel m_i + \mathcal{H}(m)$

\mathcal{CH} : a Chameleon Hash function \mathcal{H} : a Standard Crypto. Hash

k_i : secret key given to the Sanitizer needed for I-Forge

u_i : secret key that is safely discarded

Signature:

$$d_i = \begin{cases} \mathcal{CH}_{k_i}(m_i) & \text{if } m_i \text{ is sanitizable} \\ \mathcal{CH}_{u_i}(m_i) & \text{if } m_i \text{ is not sanitizable} \end{cases}$$

$$\sigma = SIGN(d_1 \parallel \dots \parallel d_n \parallel \mathcal{H}(m))$$

Mixing yields: Strong Transparency w/o Transparency

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$$\sigma = SIGN(d_1 \parallel \dots \parallel d_n \parallel \mathcal{H}(m))$$

Properties: **NO TRANSPARENCY & STRONG TRANSPARENCY**

The Property of Transparency Revisited (1)

Existing definitions of Transparency:

- T always implies WT or ST
- WT or ST always implies T
- T \Leftrightarrow ST or WT

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Practically, a verifier either knows which m_i is potentially sanitizable or he does not:

- $\models (ST \text{ or } WT)$
- **The result: T is always true**

The Property of Transparency Revisited (1)

Existing definitions of Transparency:

- T always implies WT or ST
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Practically, a verifier either knows which m_i is potentially sanitizable or he does not:

- \models (ST or WT)
- The result: T is always true

We found this counter intuitive and
the Mixing Example showed no T , but still has ST

The Property of Transparency Revisited (2)

**Transparency makes a statement about
the detection of a sanitized document.**

**Weak and Strong Transparency make statements about
the detection of sanitizable subdocuments.**

**Weak and Strong Transparency are
independent from Transparency.**

Conclusion

- 1. JAVA Implementation of 5 Sanitizable Sign. Schemes:**
 - Mostly a tolerable performance penalty over SHA/RSA
 - Full JCA integration as JAVA Crypto Provider
 - Integration into XML Digital Signature Syntax and Processing Standard (W3C)
- 2. Mixing Signature / Hash Algorithms is easy and “natural” using XML’s <references>**
 - Allows fine-grained control over Scheme’s properties like Transparency
- 3. Property of Transparency is independent from Strong-/Weak-Transparency and has a different scope**

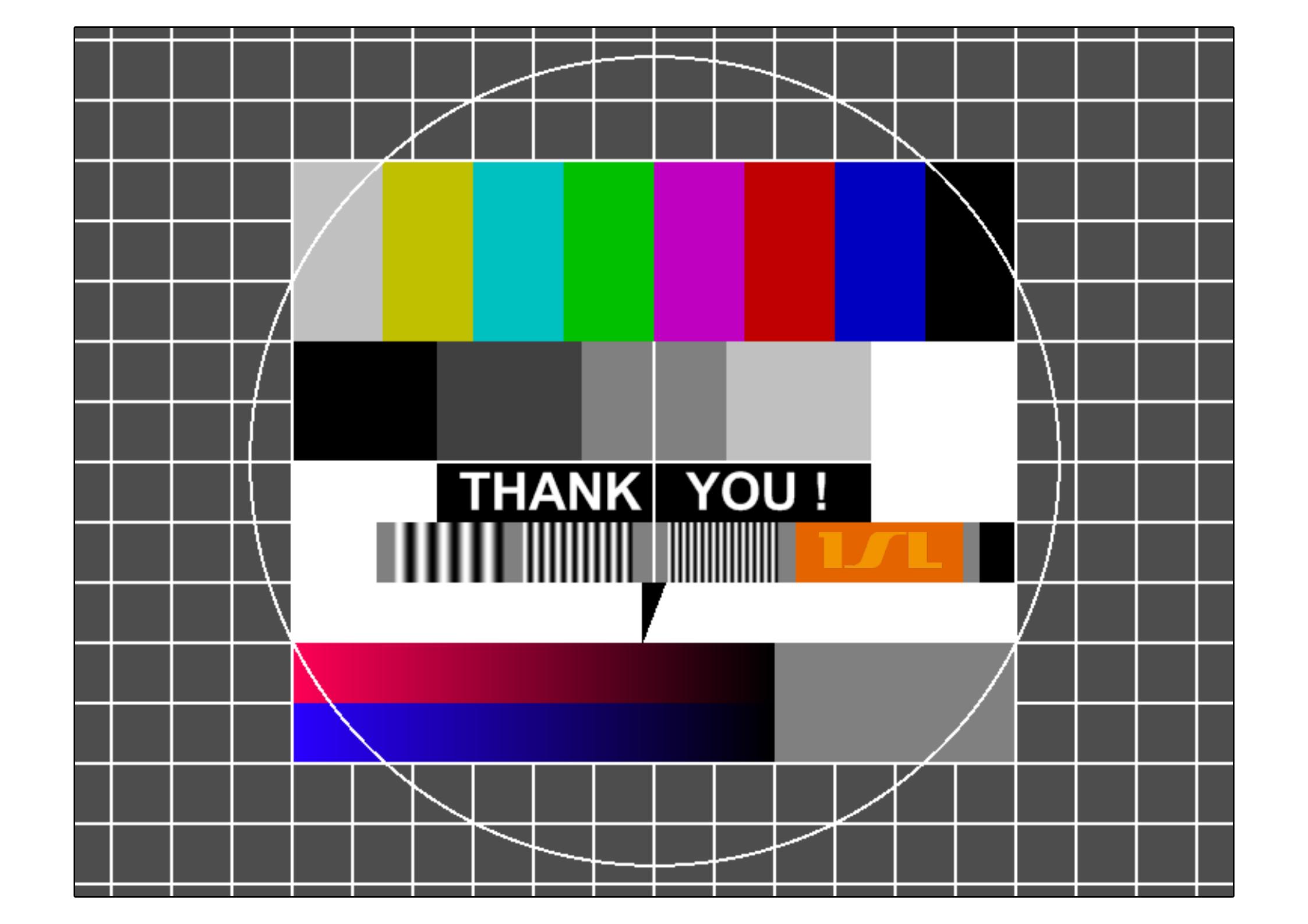
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2. Mixing Signature / Hash Algorithms is easy and “natural” using XML’s <references>

- Allows fine-grained control over Scheme’s properties like Transparency

3. Property of Transparency is independent from Strong-/Weak-Transparency and has a different scope



THANK YOU !

ISL

The Property of Transparency by Ateniese et al.

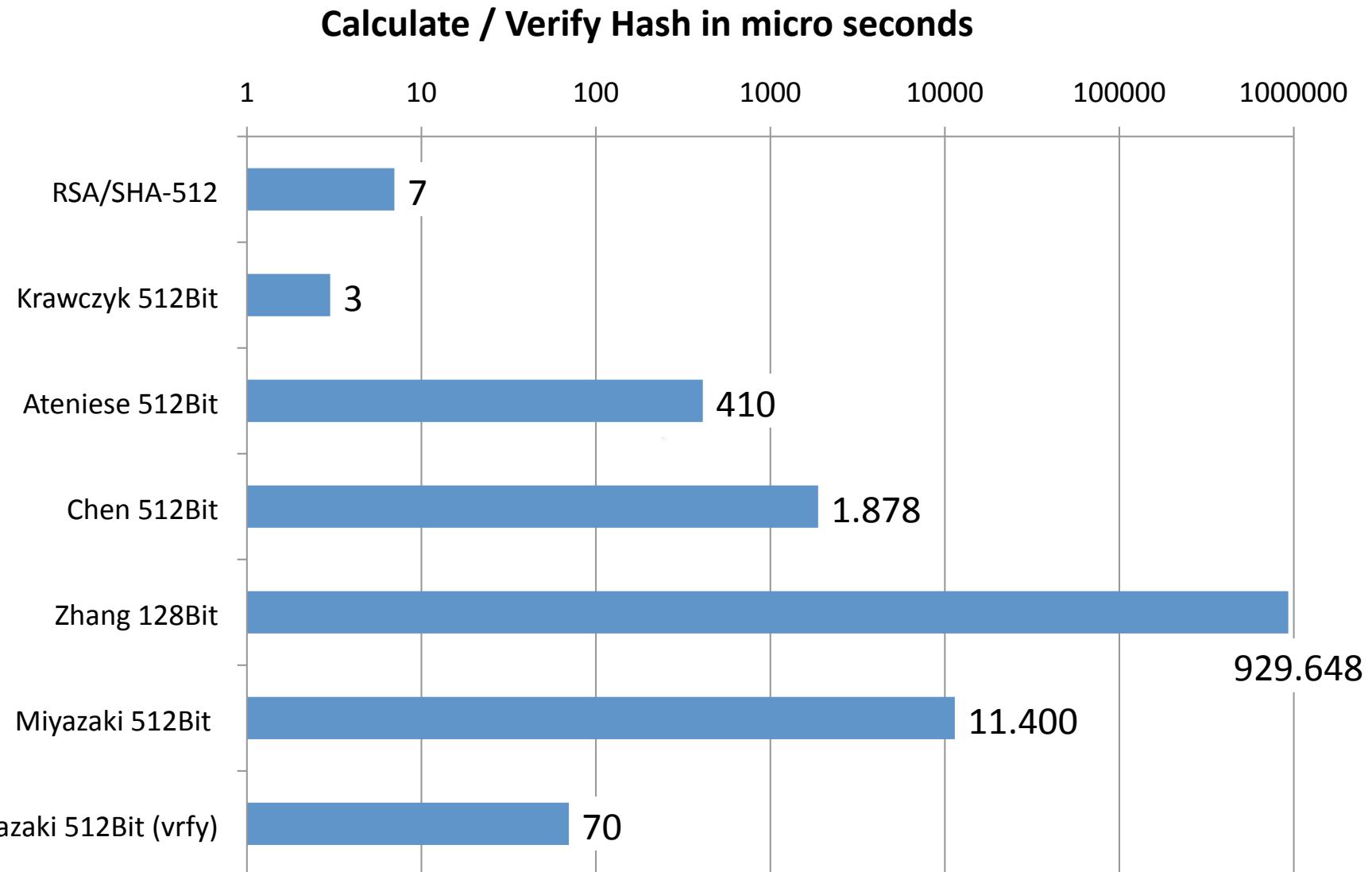
Ateniese et al. define the property of transparency (T) as follows:

Given a signed message with a valid signature, no party — except the censor and the signer — should be able to correctly guess whether the message has been sanitized. [1]

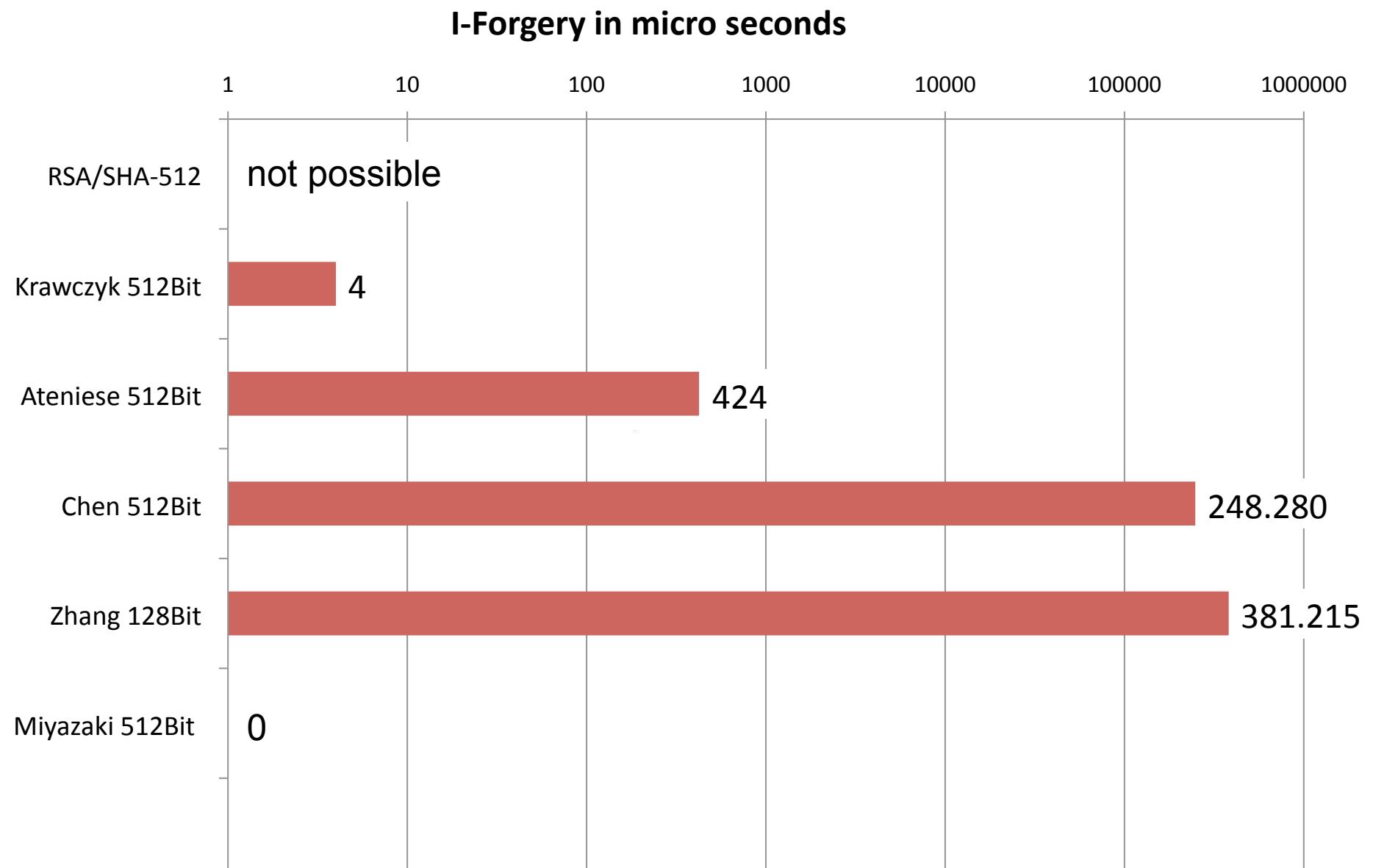
They further divide the property into “weak” (WT) and “strong transparency” (ST):

We further distinguish among two flavors of transparency: weak and strong. Weak transparency means that the verifier knows exactly which parts of the message are potentially sanitizable and, consequently, which parts are immutable. In contrast, strong transparency guarantees that the verifier does not know which parts of the message are immutable and thus does not know which parts of a signed message could potentially be sanitizable. [1]

Performance: Calculate / Verify Hash



Performance: Calc. a Forgery using I-Forge Algorithm



Performance: Details

Test Setup:

- Intel T8300 Dual Core @ 2.40 Ghz and 4 GiB of RAM.
- Algorithms coded in JAVA
- Not optimized
- make heavy use of JAVA's BigInteger class
- Integrated into JAVA Cryptographic Framework (JCA) without modifying the JCA
- Input: XML File (JAVA JCA Signature Example File)
 - 1 Reference
 - Fixed size (achievable by applying Standard Crypto. Hash 1st)